

# **Data Structures and Algorithms**

## **CS-206**

### **Queue**

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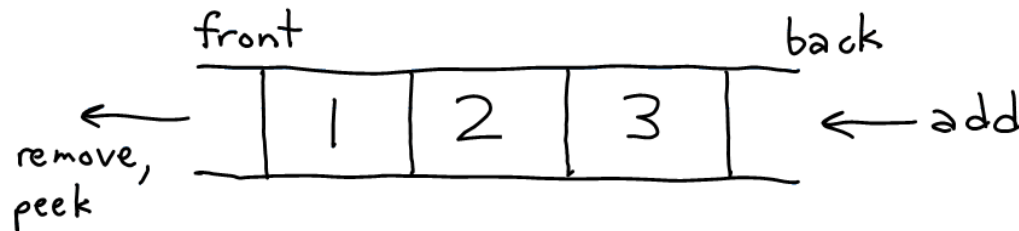


# Queue Strategy

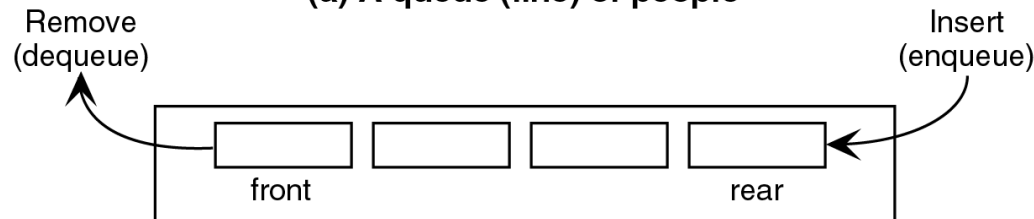
## First-in, First-out (FIFO)

Objects enter the line at one end (rear)

Objects leave the line at the other end (front)

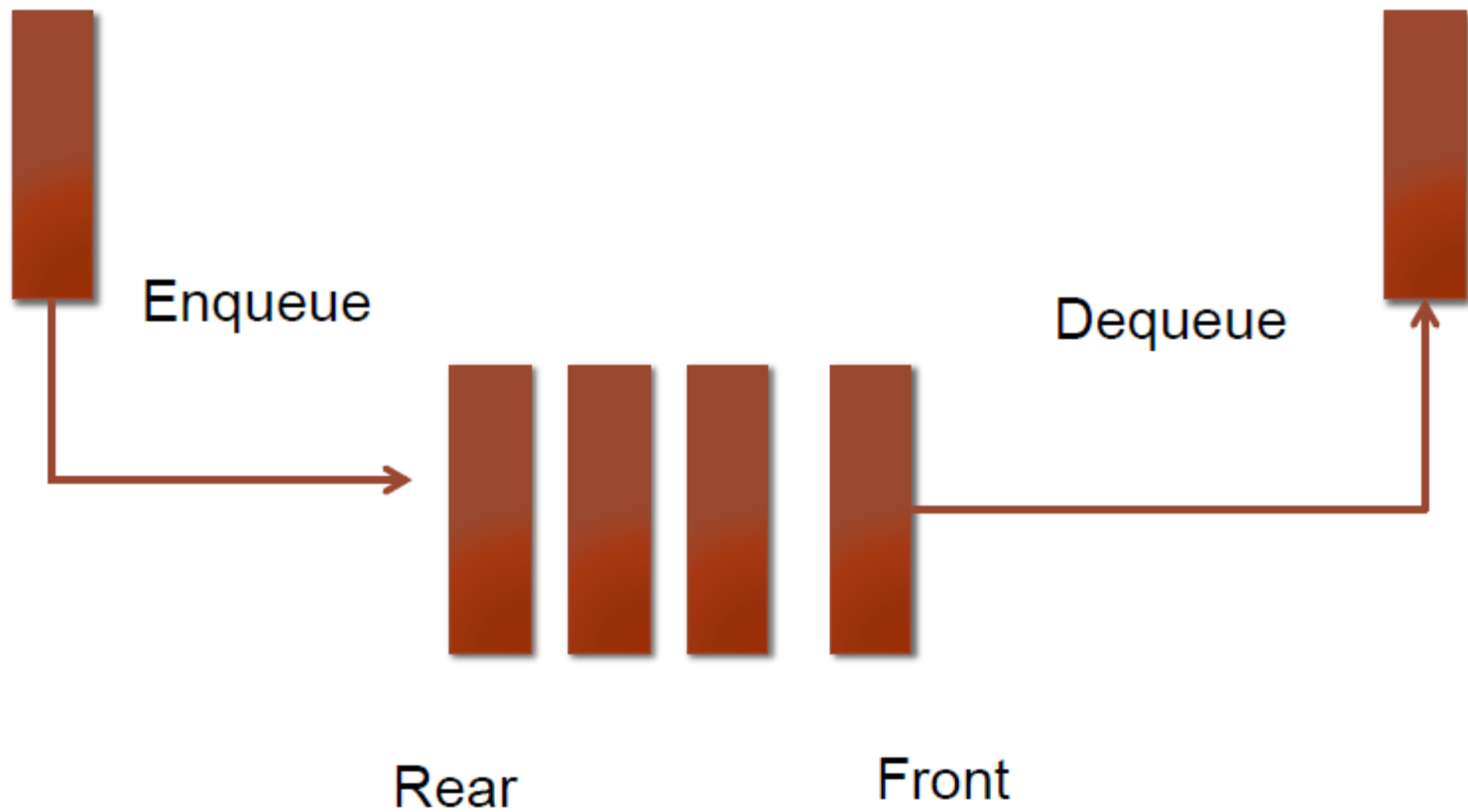


(a) A queue (line) of people



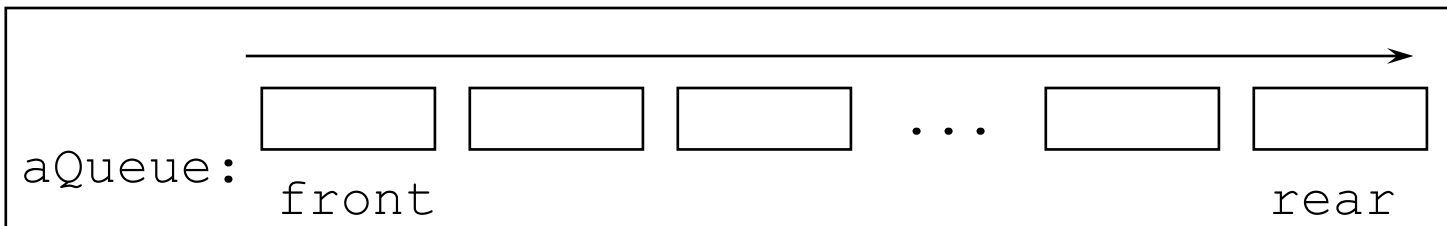
(b) A computer queue

# Queue Strategy



# Queue Definition

- Queue: Ordered collection,  
accessed only at the front (remove)  
and rear (insert)
  - Front: First element in queue
  - Rear: Last element of queue
- FIFO: First In, First Out
- Footnote: picture can be drawn in  
any direction



# Queue Example

Operation	output	queue
• enqueue(5)	-	(5)
• enqueue(3)	-	(5, 3)
• dequeue()	5	(3)
• enqueue(7)	-	(3, 7)
• dequeue()	3	(7)
• front()	7	(7)
• dequeue()	7	()
• dequeue()	"error"	()
• isEmpty()	true	()
• enqueue(9)	-	(9)
• size()	1	(9)

# Queues in Computer Science

- Operating systems:
  - queue of print jobs to send to the printer
  - queue of programs / processes to be run
  - queue of network data packets to send
- Programming:
  - modeling a line of customers or clients
  - storing a queue of computations to be performed in order
- Real world examples:
  - people on an escalator or waiting in a line
  - cars at a gas station (or on an assembly line)

# Queue Operations

- Initialize the queue,  $Q$ , to be the empty queue.
- Determine whether or not if the queue  $Q$  is empty.
- Determine whether or not if the queue  $Q$  is full.
- Insert (enqueue) a new item onto the rear of the queue  $Q$ .
- Remove (dequeue) an item from the front of  $Q$ , provided  $Q$  is nonempty.



# Queue Operations

- `enqueue()` Adds an element to rear of queue succeeds unless the queue is full (if implementation is bounded)

Or `insert()`

- `dequeue()` Remove and return the front element of queue,  
Precondition: queue is not empty

Or `remove()`

- `Front()` Return a copy of the front element of queue, precondition: queue is not empty
- `Size()`
- `isEmpty()`
- `isFull()`

# Queue Example

- Draw a picture and show the changes to the queue in the following example:
  - Queue q; Object v1, v2;
  - q.insert("chore");
  - q.insert("work");
  - q.insert("play");
  - v1 = q.remove();
  - v2 = q.front();
  - q.insert("job");
  - q.insert("fun");

# What is the result of:

- Queue q; Object v1,v2,v3,v4,v5,v6
- q.insert("Sue");
- q.insert("Sam");
- q.insert("Sarah");
- v1 = q.remove( );
- v2 = q. front( );
- q.insert("Seymour");
- v3 = q.remove( );
- v4 = q.front( );
- q.insert("Sally");
- v5 = q.remove( );
- v6 = q. front( );

# Applications: Job Scheduling

front	rear	Q[0]	Q[1]	Q[2]	Q[3]	Comments
-1	-1					queue is empty
-1	0	J1				Job 1 is added
-1	1	J1	J2			Job 2 is added
-1	2	J1	J2	J3		Job 3 is added
0	2		J2	J3		Job 1 is deleted
1	2			J3		Job 2 is deleted

# Implementing Queue ADT: Array

## Queue

- Keep track of the number of elements in the queue, `size`.
- Enqueue at the back of the array (`size`).
- Dequeue at the front of the array (index 0).
  - what is bad about this implementation?
  - what if we enqueue at 0 and dequeue at `size`?

# Array Implementation

Any implementation of a queue requires:  
storage for the data as well as  
markers (“pointers”) for the front and for the back of the queue.

An **array-based** implementation would need structures like  
**items**, an array to store the elements of the queue  
**Front**, an index to track the front queue element  
**Rear**, an index to track the position *following* last queue element

**Additions** to the queue would result in incrementing **Rear**.

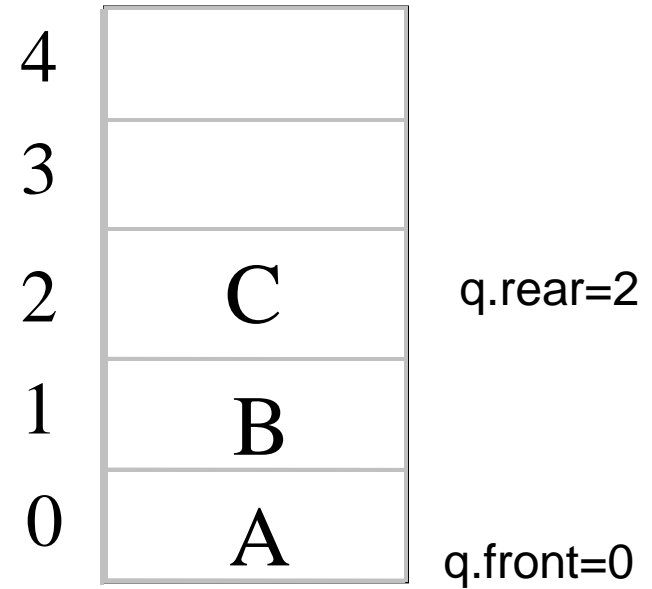
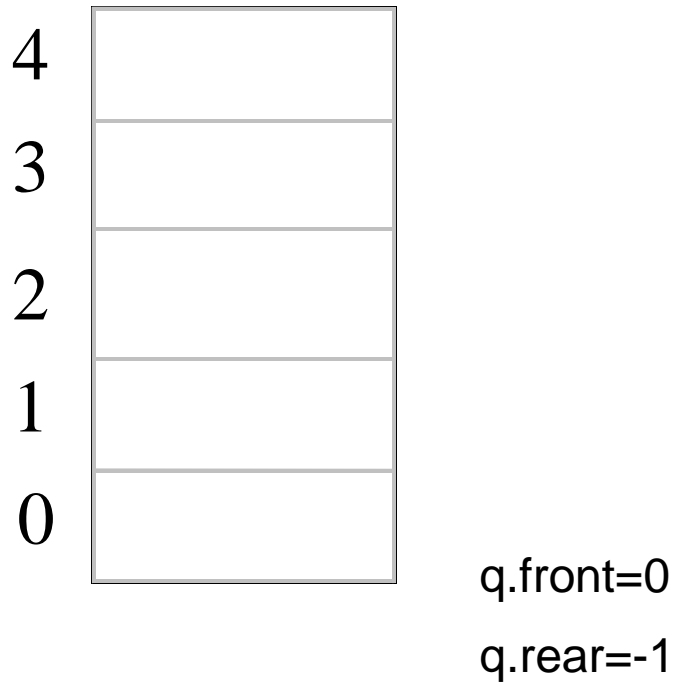
**Deletions** from the queue would result in incrementing **Front**.

*Clearly, we'd run out of space soon!*

# Queue using Arrays

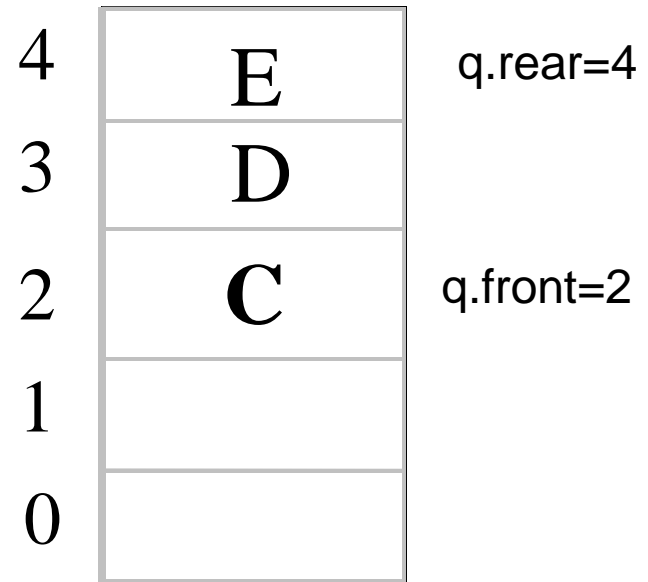
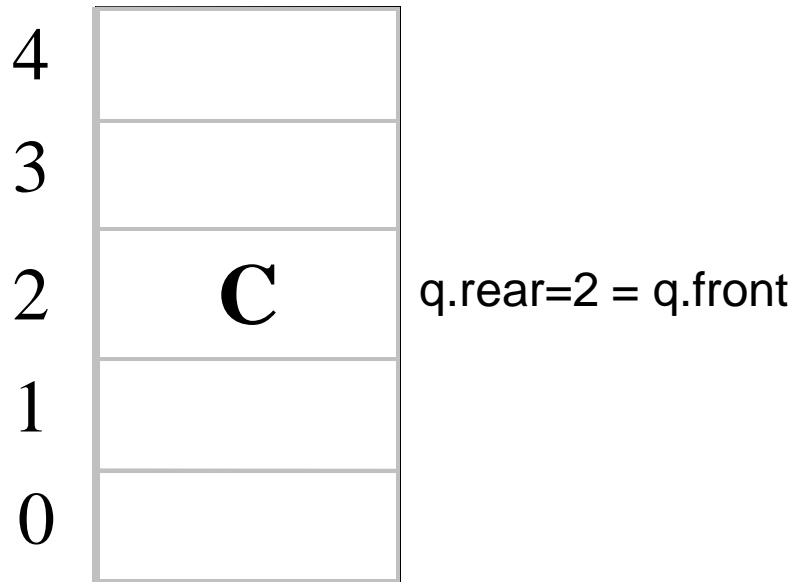
```
#define length 10
Struct queue
{
    int items[length];
    int front, rear;
}
```

- Insert(q,x)  
q.items[++q.rear]=x;
- X=remove(q)  
x=q.items[q.front++];



- The queue is empty whenever  $q.rear < q.front$
- The number of elements in the queue at any time is equal to the value of  $q.rear - q.front + 1$





- Now there are 3 elements in the queue but there is room for 5
- If to insert *F* in the queue the *q.rear* must be increased by 1 to 5 and  $q.items[5]$  must be set to the value F. but  $q.items$  is an array of only 5 elements so this insertion cannot be made

# Solution to Problem

*Clearly, we've run out of space*

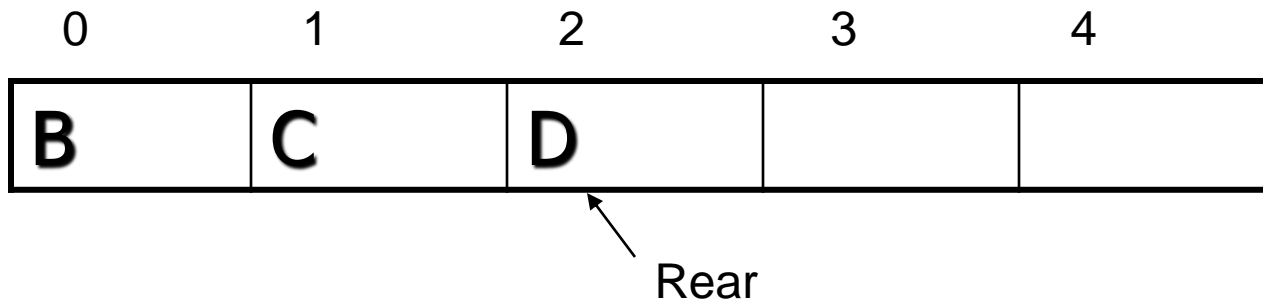
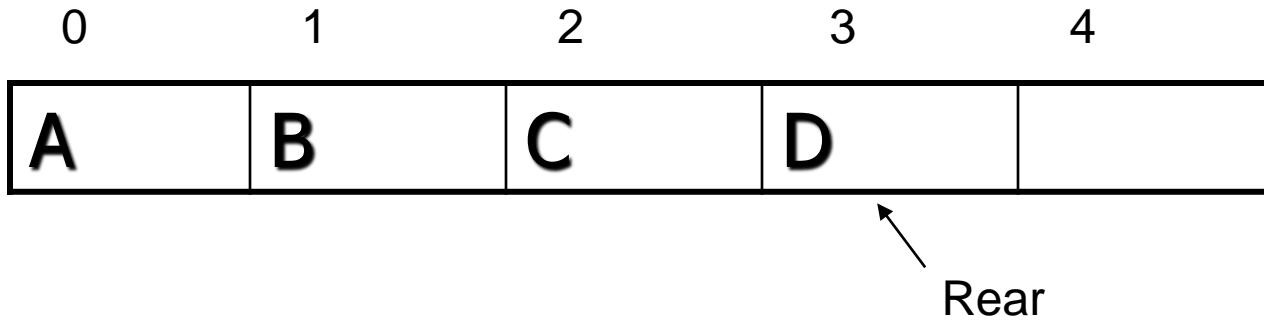
## **Solutions include:**

- Shifting the elements downward with each deletion
- Viewing array as a circular buffer, i.e. wrapping the end to the front

# Solution 1

- For arrays there are two methods
- First is do it as we do in real world
  - Check if array is not empty
  - Simply dequeue from the first location of array say `array[0]` i.e. the zero<sup>th</sup> index
  - After dequeue shift all the elements of the array from `array[index]` to `array[index-1]`

# Deque



DE-QUEUE

# Deque Operation

```
x=q.items[0];
```

```
for(i=0; i<q.rear; i++)
```

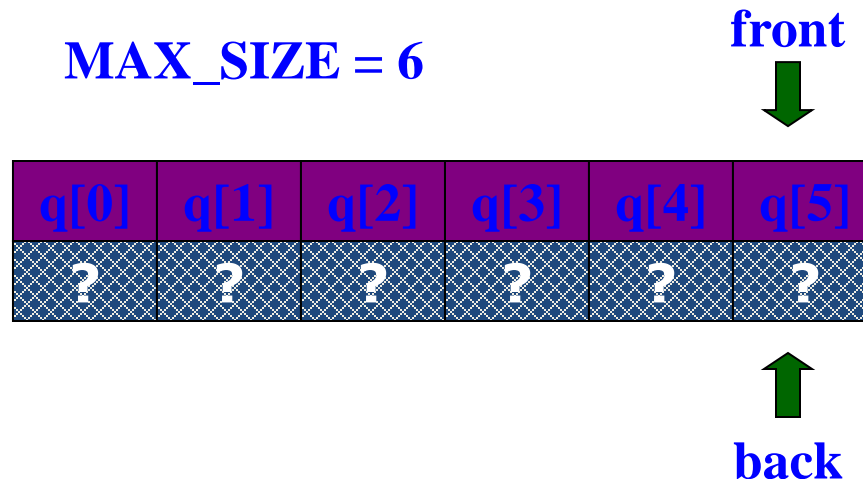
```
    q.items[i]=q.items[i+1]
```

```
q.rear--;
```

- Method is costly as we have to move all the elements of the array
- Do we need Front Index in this case?
  - No, Because we are always dequeue(ing) from the first index

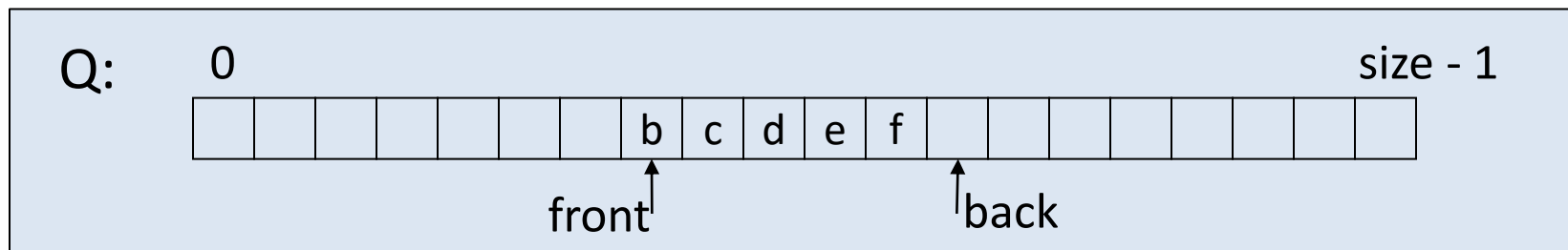
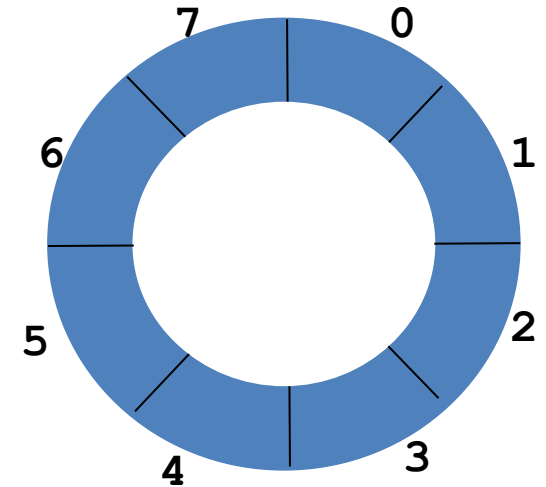
# Solution2: Circular Queue

- To avoid the costly operation of coping all the elements again we employ another method called **circular queue**
- Simply dequeue the element and move the front pointer to next index
- If  $q.rear == q.front$  queue is empty
- $q.front = q.rear = \text{max\_size} - 1$ ;



# Implementing Queue ADT: Circular Array Queue

- **Neat trick:** use a ***circular array*** to insert and remove items from a queue in constant time.
- The idea of a circular array is that the end of the array “wraps around” to the start of the array.



# Summary



- Stacks and Queues
  - Specialized list data structures for particular applications
- Stack
  - LIFO (Last in, first out)
  - Operations: push(Object), top( ), and pop( )
- Queue
  - FIFO (First in, first out)
  - Operations: insert(Object), getFront( ), and remove( )
- Implementations: arrays or lists are possibilities for each





# Assignment

Palindromes are words which can be read same from forward and revers. Few examples are:

- Radar
- Mom
- Dad
- Stats
- Madam
- Wassamassaw

How we may use Stack and Queue to determine a given word is palindrome?

# Credit

Data Structures and Algorithms in C++ Goodrich,  
Tamassia and Mount (Wiley, 2004)

University of Washington